



O N L I N E

LEARNING

Algorithm Efficiency Analysis

SCSJ2013 Data Structures & Algorithms

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Objectives

Students are expected to be able to do the following:

Understand algorithm efficiency analysis.

Able to measure algorithm efficiency using big O notation.



What is algorithm analysis?

Study the **efficiency** of algorithms when the **input size grow**, based on the **number of steps**, the amount of computer **time and the space usage**.

Analysis of algorithms

- Algorithm analysis concern with the **size and growth of data run on a particular algorithm.**
- However, algorithm analysis should be independent of :
 - Specific implementations (programming language such as C, C++ or Java)
 - Specific Computer hardware (computer chips, OS, or speed)
 - Particular set of data (string, int, float)

Analysis of algorithms

Three possible states in algorithm analysis:

Worst case

- Longest running time for *any* input of size n
- A determination of the maximum amount of time that an algorithm requires to solve problems of size n

Best Case

- Shortest running time for *any* input of size n
- A determination of the minimum amount of time that an algorithm requires to solve problems of size n

Average Case

- Average running time for *all* inputs of size n
- A determination of the average amount of time that an algorithm requires to solve problems of size n



Big O Notation

- Complexity time can be represented by **Big 'O' notation**.
- Notation that used to show the complexity time of algorithms.
- Big 'O' notation is denoted as $O(f(n))$.

whereby

O – “the order of”

$f(n)$ - algorithm's **growth-rate function**

Example, $O(1)$, $O(\log_x n)$, $O(n)$, $(O_n \log_x n)$, $O(n^2)$

Big O Notation Example

Notation	Execution time
$O(1)$	Constant function, independent of input size, n Example: Finding the first element of a list.
$O(\log_x n)$	Problem complexity increases slowly as the problem size increases. Example: Solve a problem by splitting into constant fractions of the problem (e.g., throw away $\frac{1}{2}$ at each step)
$O(n)$	Problem complexity increases linearly with the size of the input, n Example: counting the elements in a list.



Big O Notation Example cont..

$O(n \log_x n)$	<p>Log-linear increase - Problem complexity increases a little faster than n</p> <p>Characteristic: Divide problem into subproblems that are solved the same way</p> <p>Example: mergesort</p>
$O(n^2)$	<p>Quadratic increase.</p> <p>Problem complexity increases fairly fast, but still manageable</p> <p>Characteristic: Two nested loops of size n</p>
$O(n^3)$	<p>Cubic increase.</p> <p>Practical for small input size, n.</p>
$O(2^n)$	<p>Exponential increase - Increase too rapidly to be practical</p> <p>Problem complexity increases very fast</p> <p>Generally unmanageable for any meaningful n</p> <p>Example: Find all subsets of a set of n elements</p>

Big O Notation Algorithm

Notation	Codes
$O(1)$ <i>Constant</i>	<pre>int counter = 1; cout << "cout execution times" << counter;</pre>
$O(\log_x n)$ <i>Logarithmic</i>	<pre>int counter = 1; int i = 0; for (i = x; i <= n; i = i * x) { // x must be > than 1 cout << "cout execution times" << counter ; counter++; } // Ex: if x = 2 and n = 16 // i = 2, 4, 8, 16</pre>

Big O Notation

<p>$O(n)$ Linear</p>	<pre>int counter = 1; int i = 0; for (i = 1; i <= n; i++) { cout << "cout execution times" << counter ; counter++; }</pre>
<p>$O(n \log_x n)$ Linear Logarithmic</p>	<pre>int counter = 1; int i = 0; int j = 1; for (i = x; i <= n; i = i * x) { // x must be > than 1 while (j <= n) { cout << "cout execution times" << counter; counter++; j++; } }</pre>

Big O Notation

- Example of algorithm for common function:

$O(n^2)$

Quadratic

```
int counter = 1;
int i = 0;
int j = 0;
for (i = 1; i <= n; i++) {
    for (j = 1; j <= n; j++) {
        cout << "cout execution times" << counter;
        counter++;
    }
}
```

Big O Notation Algorithm

$O(n^3)$
Cubic

```
int counter = 1;
int i = 0;
int j = 0;
int k = 0;
for (i = 1; i <= n; i++) {
    for (j = 1; j <= n; j++) {
        for (j = 1; j <= n; j++) {
            cout << "cout execution times " << counter;
            counter++;
        }
    }
}
```



Big O Notation Algorithm

$O(2^n)$

Exponential

```
int counter = 1;
int i = 1;
int j = 1;
while (i <= n) {
    j = j * 2;
    i++;
}
for (i = 1; i <= j; i++) {
    cout << "cout execution times" << counter;
    counter++;
}
```



Order of increasing complexity

$$O(1) < O(\log_x n) < O(n) < O(n \log_2 n) < O(n^2) < O(n^3) < O(2^n)$$

Notasi	n = 8	n = 16	n = 32
$O(1)$	1	1	1
$O(\log_2 n)$	3	4	5
$O(n)$	8	16	32
$O(n \log_2 n)$	24	64	160
$O(n^2)$	64	256	1024
$O(n^3)$	512	4096	32768
$O(2^n)$	256	65536	4294967296



Determine the number of steps

Algorithm	No. of Steps
void sample8 ()	0
{	0
int n, x, i=1;	1
while (i<=n)	n
{	0
x++;	n.1 = n
i++;	n.1 = n
}	0
}	0
Number of Steps	1 + 3n

Consider the largest factor : 3n

and remove the coefficient : O(n)

Conclusion and Summary

Algorithm analysis to study the efficiency of algorithms when the input size grow, based on the number of steps, the amount of computer time and **space**

Can be done using Big O notation by using growth of function.

Order of growth for some common function:

• $O(1) < O(\log_x n) < O(n) < O(n \log_2 n) < O(n^2) < O(n^3) < O(2^n)$

Three possible states in algorithm analysis best case, average case and worst case.

**Thank
You**



<http://comp.utm.my/>