



# SEARCHING

# DATA STRUCTURE AND ALGORITHMS

FACULTY OF COMPUTING UNIVERSITI TEKNOLOGI MALAYSIA UTM





#### **OBJECTIVES FOR STUDENTS**

- 1. Able to describe the searching technique concept and the purpose of searching operation.
- 2. Develop C++ code to implement the basic searching algorithms.
- 3. Able to analyze the efficiency of the searching technique.
- 4. Able to develop C++ code to implement searching technique in problem solving.

# **KEY CONCEPT**

# **1.0 INTRODUCTION TO SEARCHING**

1.1.	<b>Searching definition</b> - A process to determine whether an element is a <b>member</b> of a certain data set or a collection of elements. The process also aims to find the location of the element with a specific value (key) within the collection of elements.								
1.2.	<ul> <li>The process can also be seen as an attempt to search for a certain record in a file</li> <li>i. Each record contains data field and key field</li> <li>ii. Key field is a group of characters or numbers used as an identifier for each record</li> <li>iii. Searching can done based on the key field.</li> </ul>								
1.3.	<ul> <li>.3. Consider the following data set of employee record. Searching can be done based on certain field: emplD, or empl_IC, or empName.</li> <li>To search empID = 122, give us the record value at index 1.</li> </ul>								
inde	ex	empID	Empl_IC	EmpName	Post				
	<b>ex</b> [0]								
	-	emplD	Empl_IC	EmpName	Post				
	[0]	<b>emplD</b>	Empl_IC 701111-11-1234	EmpName Ahmad Faiz Azhar Mohd. Azim Bin	<b>Post</b> Programmer				





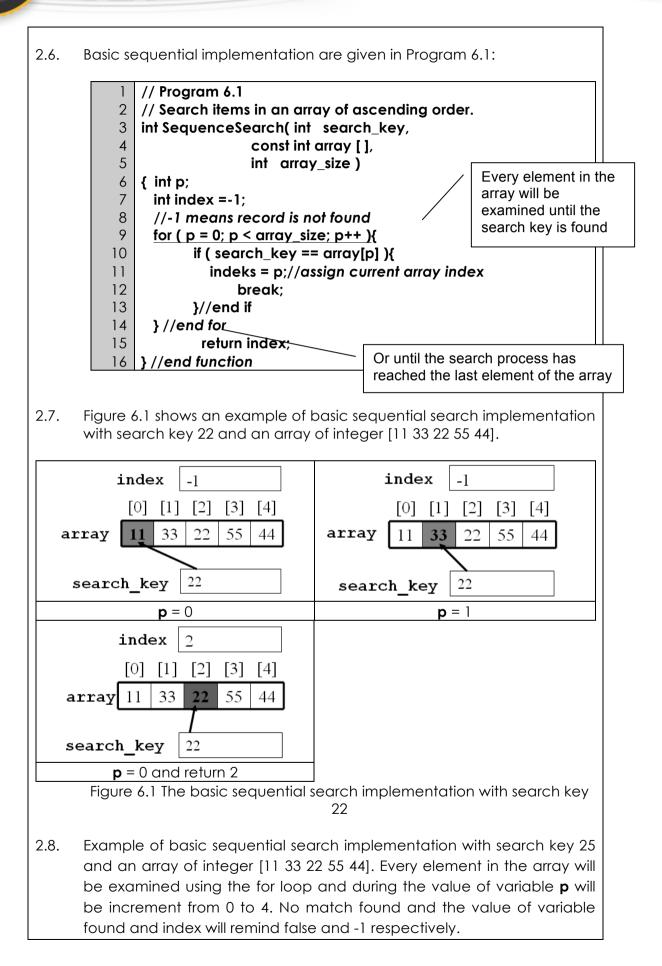
- i. Sequential search
- ii. Binary Search
- iii. Binary Tree Search
- iv. Indexing
- 1.5. Similar with sorting, Searching can also be implemented in two cases, **internal** and **external** search.
  - External search only implemented if searching is done on a very large size of data. Half of the data need to be processed in RAM while half of the data is in the secondary storage.
  - Internal search searching technique that is implemented on a small size of data. All data can be load into RAM while the searching process is conducted.

# 2.0 BASIC SEQUENTIAL SEARCH

- 2.1. Used for searching that involves records stored in the main memory (RAM).
- 2.2. Basic sequential search also used to search an element from unsorted list.
- 2.3. Basic sequential search is the **simplest search algorithm**, but is also the slowest and can only be used to search from a small list. The efficiency of sequential search is low compared to other searching techniques.
- 2.4. In a sequential search, (1) every element in the array will be examine sequentially, starting from the first element; (2) The process will be repeated until the last element of the array or until the searched data is found.
- 2.5. Searching strategy :
  - **Examines** each element in the array one by one (sequentially) and compares its value with the one being looked for the search key.
  - Search is successful if the search key **matches** with the value being compared in the array. Searching process is **terminated**.
  - else, if no matches is found, the search process is continued to the last element of the array. Search is failed array if there is no match found from the array.



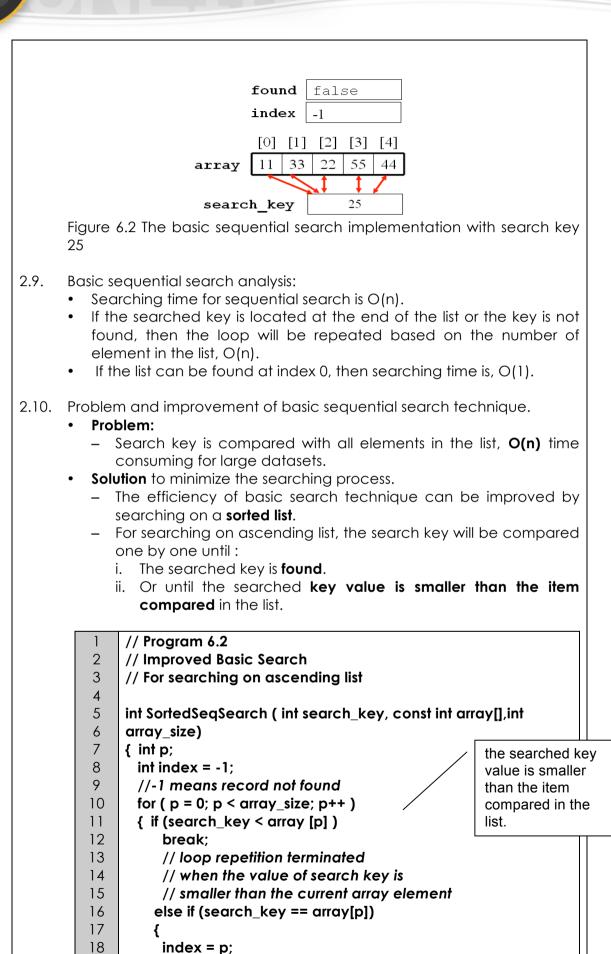














the searched

key is found

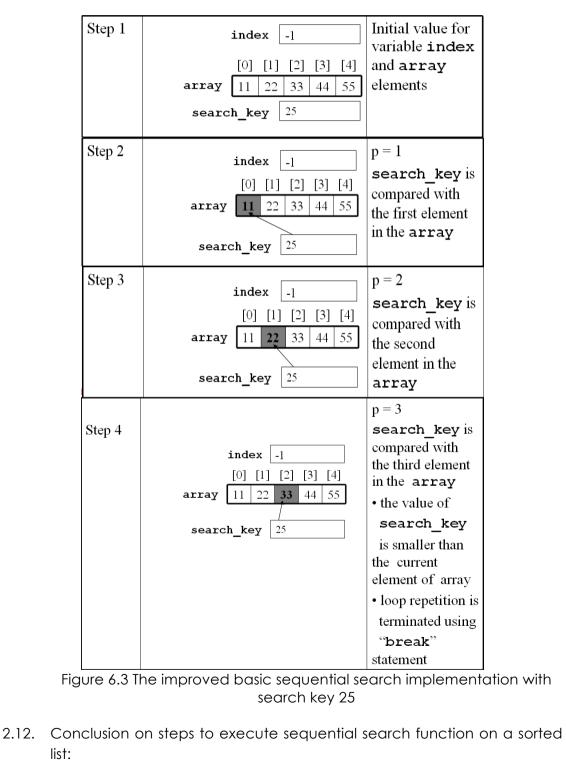




20 break; 21 }// end

19

- } // end else-if
- 22 }//end for
- 23 return index; // return the value of index
- 24 } //end function
- 2.11. Example of the improved basic sequential search with search key 25 and an array of integer [11 22 33 44 55]:





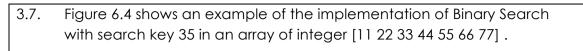


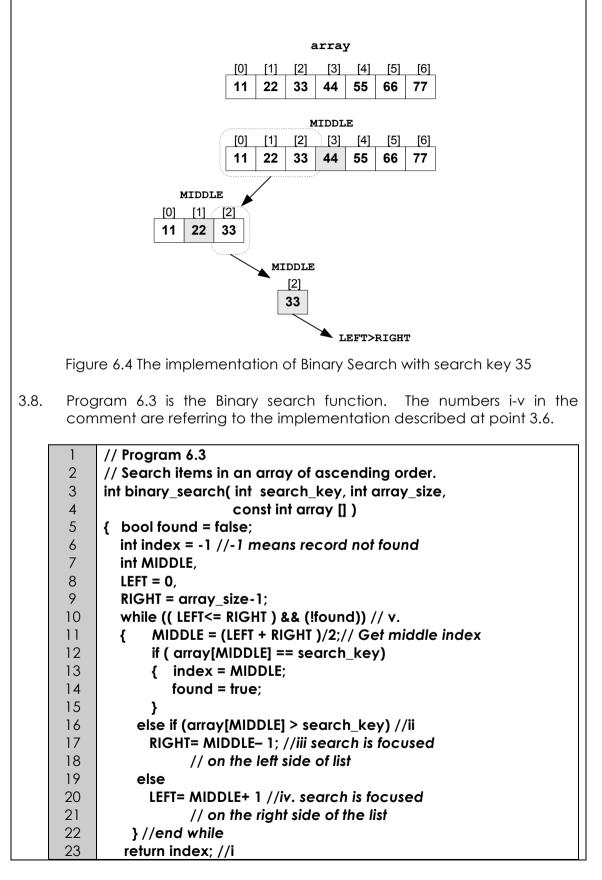
- If the elements in the list are not in a sorted (asc/desc) order, loop will be repeated based on the number of elements in the list.
- When the list is not sorted the loop is repeated 5 times, compared to 3 times if the list is in sorted order as shown in the previous example.
- If the list is sorted in descending order, change operator "<" to operator ">" in the loop for.

# **3.0 BINARY SEARCH**

- 3.1. The drawback of sequential search algorithm is having to **traverse the** entire list, O(n).
- 3.2. Sorting the list does minimize the cost of traversing the whole data set, but we can improve the searching efficiency by using the **Binary Search** algorithm.
- 3.3. Consider a list in ascending sorted order. For a sequential search, searching is from the beginning until an item is found or the end is reached.
- 3.4. Binary search improve the algorithm by **removing as much of the data** set as possible so that the item is found more quickly.
- 3.5. Search process is started at the **middle** of the list, then the algorithm determine which half the item is in (because the list is sorted).
  - It divides the working range in half with a single test. By repeating the procedure, the result is an efficient search algorithm-O(log2 n).
- 3.6. Implementation of Binary Search starts by comparing the search key with the element at the **middle**.
  - i. If the value matches, it will be return to the calling function (index = MIDDLE)
  - ii. If the search key < the middle element, search will be focused on the elements between the first element to the element before the middle element (MIDDLE -1)
  - iii. If the search key is not found, the element at the middle of the first element and the MIDDLE -1 element will be compared with the search key.
  - iv. If the search key > the middle element, search will only be focused on the elements between the second MIDDLE element to the first MIDDLE element.
  - v. Search is repeated until the searched key is found or the last element in the subset is traversed (LEFT > RIGHT).













#### 24 }//end function

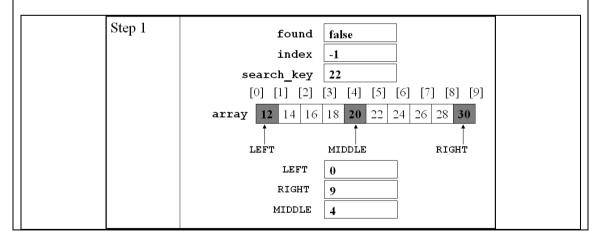
- 3.9. Consider the implementation of the Binary Search on a sorted list [11 22 33 44 55 66 77] with search key 35.
  - a. Search starts by obtaining the **MIDDLE** index of the array: **MIDDLE**= (0+6) / 2= 3 { First **MIDDLE** index}
  - b. search\_key **35** is compared with the element at the fourth index in the array, which is array[3] with the value **44**.
  - c. search\_key < MIDDLE value, therefore search will be focused on the elements between the first index and the third index only (index 1 to MIDDLE-1)
  - d. Process to obtain MIDDLE index is repeated:
     MIDDLE = (0 + 2) / 2= 1 { second MIDDLE index}
  - e. search\_key 35 is compared with the element at the second index, which is array[1] with the value 22
  - f. search\_key > MIDDLE value, therefore search will be focused only on the elements between the second MIDDLE index to the first MIDDLE index.

**MIDDLE** = (2 + 2) / 2= 2 { third **MIDDLE** index}

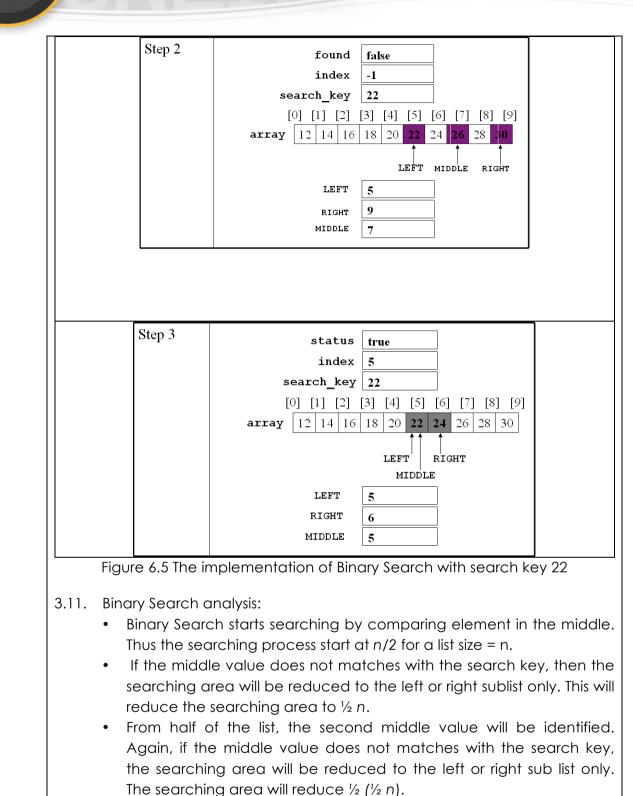
- g. Element at the third index, array[2] with the value 33 is not equal to the value of the search key.
- h. Search process has reached the last element of the traversed subset, therefore search is terminated and assumed fail.
- To search from the list sorted descending, change operator " > " to operator " < " to the following statement :</li>

else if (array [ MIDDLE ] > search\_key) RIGHT = MIDDLE – 1;

3.10. Consider another implementation of the Binary Search on a sorted list [12 14 16 18 20 2 24 26 28 30] with search key 22 and the array size is 10 in Figure 6.5.







- The process of looking for middle point and reducing the searching area to the left or right sublist will be repeated until the middle value is equal to the middle value (search key is found) or the last value in sublist has been traverse.
- If the repetition occur k times, then at iteration k, the searching area is reduced to  $(\frac{1}{2})^k n$ .
- Figure 6.6 shows the reducing size for binary search area.





	<ul> <li>At iteration k for array size n to (<sup>1</sup>/<sub>2</sub>)<sup>k</sup>n.</li> </ul>	e = n , searching area will be reduced from								
		_								
		n								
ſ		·								
k≺										
	Figure 6.6 The reducin	ng size for binary search area								
		ig size for bindry sedicit dred								
	Eiguro 6.7 shows an ovar	and of the reducing size for hings, sag								
	_	nple of the reducing size for binary sea								
		area with 1 billion data.								
		It can be concluded that to search item in the middle of the list, the								
	complexity is O(1).									
	Searching item at the back or front of the list is faster with only $O(29)$ .									
	_									
	_									
	_									
	However, searching item									
	However, searching item	at front of the list is the worst case of Bin								
	<ul> <li>However, searching item of Search with O(29).</li> </ul>									
	<ul> <li>However, searching item of Search with O(29).</li> <li>500,000,000</li> </ul>	at front of the list is the worst case of Bin 16 15,258 17 7,629 18 3,814								
	<ul> <li>However, searching item of Search with O(29).</li> <li>500,000,000</li> <li>250,000,000</li> </ul>	at front of the list is the worst case of Bin 16 15,258 17 7,629 18 3,814 19 1,907								
	<ul> <li>However, searching item of Search with O(29).</li> <li>500,000,000</li> <li>250,000,000</li> <li>125,000,000</li> </ul>	at front of the list is the worst case of Bin 16 15,258 17 7,629 18 3,814 19 1,907 20 953								
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	<ul> <li>However, searching item of Search with O(29).</li> <li>500,000,000</li> <li>250,000,000</li> <li>125,000,000</li> <li>31,250,000</li> <li>31,250,000</li> <li>531,250,000</li> <li>615,625,000</li> <li>7,812,500</li> <li>83,906,250</li> </ul>	at front of the list is the worst case of Bin 16 15,258 17 7,629 18 3,814 19 1,907 20 953 21 476 22 238 23 119 24 59 25 29								
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	<ul> <li>However, searching item of Search with O(29).</li> <li>500,000,000</li> <li>250,000,000</li> <li>125,000,000</li> <li>31,25,000</li> <li>615,625,000</li> <li>7,812,500</li> <li>3,906,250</li> <li>1,953,125</li> <li>976,562</li> <li>488,281</li> <li>244,140</li> </ul>	at front of the list is the worst case of Bin 16 15,258 17 7,629 18 3,814 19 1,907 20 953 21 476 22 238 23 119 24 59 25 29 26 14 27 7 28 3								





# **PROGRAMMING EXERCISES**

# LAB 1: BASIC SEQUENTIAL SEARCH

Given the following Program 6.4, type and run the program to perform the tasks given below.

1	// Program 6.4
2	#include <iostream.h></iostream.h>
3	int SequentialSearch(int [], int, int);
4	
5	void main()
6	{
7	int k, target, j;
8	char ans = 'y';
9	
10	cout << "Please enter size of the array:";
11	cin >> k;
12	for(int i = 0; i < k; i++)
13	{
14	cin >> num[i];
15	}
16	do {
17	cout << "\nEnter the search key :";
18	cin >> target;
19	j = SequentialSearch(num, k, target);
20	if $(j = -1)$
21	cout << "Failed" << endl;
22	else
23	cout << "Found at num[" << j << "]\n";
24	cout << "Find another number?:";
25	cin >> ans;
26	} while (ans == 'y');
27	
28	}
29	int SequentialSearch(int a[], int n, int target)
30	{ int i;
31	for (i = 0; i < n; i++)
32	if (a[i] == target)
33	return i;
34	return -1;
35	}

Read the input for **num** array which has the following 10 ascending numbers: i.

	[0]	[1]	[2]	[3]	[4]	[5]	[6]	[7]	[8]	[9]
num	8	4	10	5	20	4	15	23	12	11
	Figure 6.8 Unsorted data									

Figure 6.8 Unsorted data





- ii. Perform searching with the following key values: 5, 4, and 25.
- iii. What are the output be when performing a search for the values?

#### LAB 2: BINARY SEARCH

Refer to the given Program 6.4 in Lab 1, replace the **SequentialSearch()** function with **BinarySearch()**function given below.

1	// Program 6.5
2	int BinarySearch(int a[], int n, int target)
3	{ int first = 0;
4	int last = n - 1;
5	int mid;
6	while (first <= last)
7	{ mid = (first + last) / 2;
8	if(target == a[mid])
9	return mid;
10	else if(target < a[mid])
11	{
12	last = mid - 1;
13	cout << "Middle value:" << mid <<
14	"\tfirst:" << first << "\tlast:"
15	<< last << endl;
16	}
17	else
18	{
19	first = mid + 1;
20	cout << "Middle value:" << mid <<
21	"\tfirst:" << first << "\tlast:"
22	<< last << endl;
23	}
24	}
25	return -1;
26	}

Based on Program 6.4 and the function in Program 6.5, perform the following tasks:

a.

Read **num** array input with the following 10 ascending numbers:

	[0]	[1]	[2]	[3]	[4]	[5]	[6]	[7]	[8]	[9]
num	2	5	6	8	10	12	15	16	20	21

Figure 6.9 Sorted data for **num** array

b. Perform search with the value of search key 5, 20 and 25. What will be the output?

iv.



# LAB 3: BASIC SEQUENTIAL SEARCH AND BINARY SEARCH ALGORITHMS

Carefully study the following program template and the sequential searching function on a sorted list called **SeqSearch**().

1	//Program 6.6
2	#include <iostream.h></iostream.h>
3	#include <conio.h></conio.h>
4	#include <stdio.h></stdio.h>
5	
6	class Person
7	{ public:
8	int key;
9	char name[30];
10	char uta[30];
11	public:
12	Person()
13	{ Person(0,'''','''');
14	}
15	Person( int key, char name[], char uta[])
16	{ this->key=key;
17	strcpy(this->name, name);
18	strcpy(this->uta, uta);
19	}
20	};
21	/********
22 23	/*************Function prototypes *******************/ void display_array(Person list[], int size, char title[]);
23 24	void QuickSort(Person list[], int first, int last);
25	int Divider(Person T[], int awal, int last);
26	int SeqSearch(int key, Person list[], int size);
27	int Binary_Search(int key, Person list[], int size);
28	void pause();
29	/**************************************
30	const int COUNT = 14; // number of array elements.
31	void main(void)
32	{ // array T, object Person
33	Person T[COUNT] = { Person(21, "Utada", "ichi"),
34	Person(61, "Hikaru", "ni"),
35	Person(11, "Ito", "san"),
36	Person(31, "Yuna", "shi"),
37	Person(79, "Hamasaki", "yon" ),
38	Person(83, "Ayumi", "go" ),
39	Person(68, "Koda", "roku" ),
40	Person(78, "Kumi", "shichi" ),
41	Person(96, "Namie", "nana" ),
42	Person(87, "Amuro", "hachi" ),
43	Person(57, "Otsuka", "ku" ),
34 25	Person(88, "Ai", "kyu"), Bergon(49, "Kaoru", "buuu")
35	Person(69, "Kaoru", "jyuu" ),

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**}**;

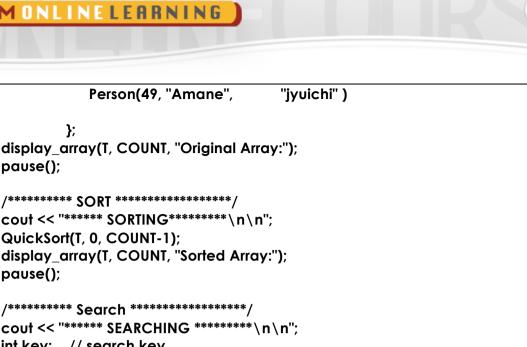
pause();



36

37 38

39 40



```
41
     /*******************************/
42
43
     cout << "****** SORTING*******\n\n";
44
     QuickSort(T, 0, COUNT-1);
     display_array(T, COUNT, "Sorted Array:");
45
46
     pause();
47
     /*********** Search **************/
48
49
     cout << "****** SEARCHING ********\n\n";
50
     int key; // search key
51
      int index; //the array index for the element found
52
            //from searching process
53
54
     cout <<"Enter search key: ";
55
     cin >> key;
56
57
     // call the search function here
58
                index = SeqSearch(key,T,COUNT);
59
     // index = Binary_Search(key,T,COUNT);
60
61
     cout << "\n\nSearch result:\n";</pre>
     cout << "\tIndex element: " << index << endl;
62
63
     if (index>-1)
     { cout << "\tKey: " << T[index].key << endl;
64
      cout << "\tName: " << T[index].name << endl;</pre>
65
      cout << "\tTitle: " << T[index].uta << endl;
66
67
     }
68
     else
69
       { cout << "\tRecord not found!!!" << endl;
70
       }
71
     pause();
72
   }
73
    74
75
     Function: display_array
76
     Description: Print object Person from array
77
    78
79
    void display_array(Person list[], int size, char title[])
80
   {int i;
81
    cout << title <<endl;
    cout <<" Key \t Name \t Title \n";
82
83
    cout <<"----- \t ------ \n";
84
85
    for (i=0; i<size; i++)
```



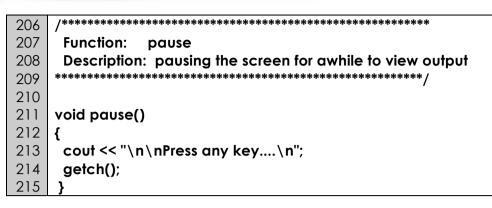


86 cout << list[i].key <<"\t"<< list[i].name 87 <<"\t\t"<< list[i].uta <<"\n"; 88 cout << endl; 89 } 90 91 92 Function: QuickSort 93 **Description: Execute Quick Sort Algorithm** \*\*\*\*\*\*\*\*\*/ 94 95 void QuickSort(Person list[], int first, int last) 96 { int cut; 97 98 if (first<last) 99 { cut = Divider(list, first, last); 100 QuickSort(list, first, cut); 101 QuickSort(list, cut+1, last); 102 } 103 } 104 105 int Divider(Person T[], int first, int last) 106 int pivot; { 107 int loop, divide, frombottom, fromtop; 108 Person temp; 109 110 pivot = T[first].key; 111 frombottom = first; fromtop = last; 112 loop = 1;113 114 while (loop) 115 { while (T[fromtop].key>pivot) 116 { // search for values less than pivot 117 // from the top of array 118 fromtop--; 119 } 120 121 while (T[frombottom].key<pivot) 122 { // search for values bigger than pivot 123 // from bottom of array 124 frombottom++; 125 } 126 127 if (frombottom < fromtop) { // swap location of pivot 128 129 temp = T[frombottom]; 130 T[frombottom] = T[fromtop]; 131 T[fromtop] = temp; 132 } else 133 { loop = 0; 134 divide = fromtop; 135



136 } //end while (loop) return divide; 137 138 } //end function Divider 139 140 151 Function: SeqSearch **Description: Execute Sequential Search Process** 152 153 Parameters: 154 key, which is the search key 155 list, which is the array, assume the array list is 156 sorted in ascending order 157 size, the number of records in array 158 **Returns:** 159 If record is found, the record index will be returned. Otherwise, if not found, value -1 will be returned. 160 161 162 163 164 int SeqSearch(int key, Person list[], int size) 165 { 166 int i: 167 for (i = 0; i < size; i++)if (key == list[i].key) return i; 168 { // if record is found 169 170 if (key < list[i].key) return -1; 181 // if search key > current record key, 182 // no need to continue search on 183 // the remaining records 184 } 185 return -1; 186 } /\*\*\*\*\*\*\*\* 187 188 **Function: Binary Search** 189 **Description: Execute Binary Search Process** 190 **Parameters:** 191 key, which is the search key 192 list, which is the array, assume the array list 193 is sorted in ascending order 194 size, the number of records in array 195 **Returns:** 196 If record is found, the record index will be returned. 197 Otherwise, if not found, value -1 will be returned. 198 \*\*\*\*\*/ 199 int Binary\_Search(int key, Person list[], int size) 200 // local variable { bool found = false: 201 202 int MIDDLE, LEFT = 0, RIGHT = size - 1; int i = -1; 203 204 } 205





Based on the given code modify the program according to the following specification:

- a. Write a new search function **Binary\_Search**() that executes binary search algorithm to replace function **SeqSearch**().
- b. Change the order of the array from ascending to descending by modifying the sort function **QuickSort**().Write again both search function **SeqSearch**() and **Binary\_Search**() based on descending order.
- c. Execute the program and see if your searching functions give out the correct output.

#### EXERCISES

#### EXERCISE 1: SEQUENTIAL SEARCH ALGORITHM

Program 6.7 is a sequential search function.

1	// Program 6.7
2	int SequenceSearch(int search_key, int array[],int array_size )
3	{ int p;
4	int index = -1;
5	// -1 means record is not found
6	
7	for (p = 0; p < array_size; p++)
8	{
9	if (search_key == array[p])
10	{ indeks = p;
11	<pre>// assign current array index</pre>
12	break; // terminate searching
13	} // end if
14	} // end for
15	return index;
16	// return location of value
17	} // end function
18	
19	

In the following figure, **DATA** is an array with the size of 10, which stores integer values.

index	[0]	[1]	[2]	[3]	[4]	[5]	[6]	[7]	[8]	[9]
DATA	49	21	34	35	40	7	15	26	12	4

Figure 6.10 Unsorted data	
---------------------------	--

- a. Based on the given function, show each steps of the search process on the **DATA** array based on the following search key values.
  - i. 49 ii. 4
  - iii. 37
- b. Compare the searching time for both search keys. Based on the searching time, discuss the efficiency of the searching technique.
- v.
- c. Explain the use of **break** statement to the **for** loop in the function in Program 6.7.

In the following figure, the array **DATA1** has been sorted in ascending order.

index	[0]	[1]	[2]	[3]	[4]	[5]	[6]	[7]	[8]	[9]
DATA1	4	7	12	15	21	26	34	35	40	49
•			F	igure	6.11 S	orted	data			<u> </u>

- d. Rewrite the function in Program 6.7 to execute searching on a sorted list.
- e. Based on your answer in question (d), show each steps of the search process on the **DATA1** array based on the following search key values.
  - i. 4 ii. 37
- f. Based on your answer in question (e), explain how and why is the sequential search on a sorted list better than on an unsorted list if the value being searched is not in the list.

# EXERCISE 2: SEQUENTIAL SEARCH ALGORITHM AND THE ANALYSIS

Assume that **r** is an array consisting *n* number of records. Each record has a key field k, and the key for the **i**-th record is referred to as **r[i].k**. The records in the array **r** are in an unsorted order.





- a. Write a sequential search function in C++, for the array **r**, with the purpose of searching a given record key. Use the most efficient algorithm that you have learned.
- b. If the array **r** are in a sorted order, write a sequential search function in C++ that will execute a search on a sorted array based on a given search key.
- c. Discuss the difference in the efficiency of the algorithm for the answers gave in questions (a) and (b).

# EXERCISE 3: BINARY SEARCH ALGORITHM

Program 6.8 is a binary search function.

1	// Program 6.8
2	int binary_search( int search_key, int array_size,
3	const int array [] )
4	{
5	bool found = false;
6	int index = -1 //-1 means record not found
7	int MIDDLE,
8	LEFT = 0,
9	RIGHT = array_size - 1;
10	
11	while (( LEFT <= RIGHT ) && (!found))
12	{ MIDDLE = (LEFT + RIGHT )/ 2;
13	// Get middle index
14	if (array[MIDDLE] == search_key)
15	{ index = MIDDLE;
16	found = true;
17	}
18	else if (array[MIDDLE] > search_key)
19	RIGHT = MIDDLE – 1;
20	<pre>// search is focused</pre>
21	<pre>// on the left side of list</pre>
22	else
23	LEFT = MIDDLE + 1
24	<pre>// search is focused</pre>
25	<pre>// on the right side of the list</pre>
26	} // end while
27	return index;
28	} // end function
29	

- a. Discuss the differences between sequential search and binary search algorithms.
- b. Based on the function in Program 6.8 show each steps of the search process on the **DATA** array in Figure 6.11 on the following search key values.





- i. 4 ii. 37
- c. State the values for variables LEFT, RIGHT and MIDDLE found at each step.

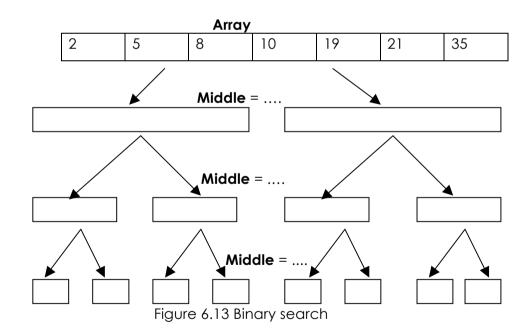
#### EXERCISE 4: BINARY SEARCH ALGORITHM

Give 2 classes of search. State which search is suitable for huge records and which is suitable for small number of records.

a. Given an ascending order array as follows:

index	[0]	[1]	[2]	[3]	[4]	[5]	[6]			
Array	2	5	8	8 10		19 21				
Figure 6.12 Sorted data										

b. Give right and left values in a box that are involved in the process of binary search. Assume that the search key is 15. Fill in the values in Figure 6.13.



#### EXERCISE 5: BASIC SEQUENTIAL SEARCH AND BINARY SEARCH ALGORITHMS

The following class diagram named **month** has 2 attributes, **key** – an integer value for **month** and **monthName** – the **char**[] value for month. The following array figure shows an array of class month, named **arrayA** with 7 elements. The array is sorted in descending order.

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	Month	
ſ	key : int	
	monthName : char[]	
		-

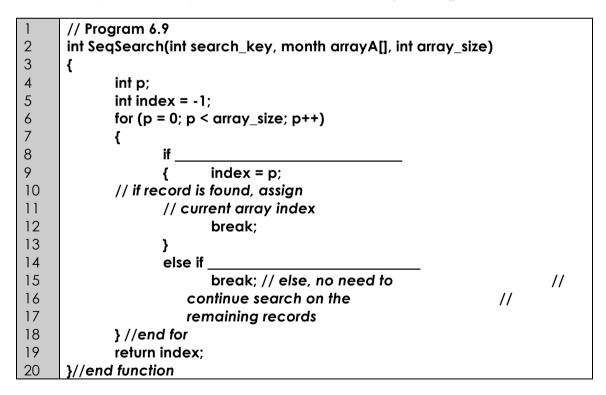
Figure 6.14 Class diagram month

	[0]	[1]	[2]	[3]	[4]	[5]	[6]
Key	12	10	8	7	5	4	2
monthNa	Decemb	Octobe	August	July	Ma	April	February
me	er	r			у		

Figure 6.15 arrayA

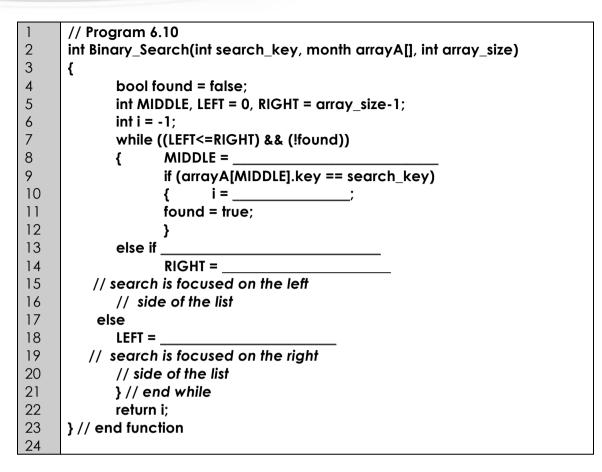
Answer the following questions based on arrayA.

a. Complete the sequeantial search function SeqSearch() below :



b. Complete the binary search function of **Binary\_Search()** below.





c. Complete the table below to show the values of the variables i, left, right, middle and found in order to search 12 as the search key during the searching process of **Binary\_Search()** function. The searching process is implemented in the **arrayA**, which has been sorted in descending order.

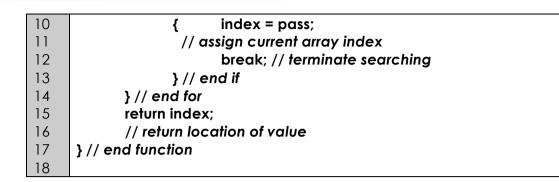
i	LEFT	RIGHT	MIDDLE	FOUND

# EXERCISE 6: BASIC SEQUENTIAL SEARCH AND BINARY SEARCH ALGORITHMS

Program 6.11 is a sequential search function. Based on the array **DATA** figure below, answer all the following questions.

1	// Program 6.11
2	<pre>int SequenceSearch(int search_key, int array[], int array_size )</pre>
3	{ int pass;
4	int index = -1;
5	//-1 means record is not found
6	
7	for (pass = 0; pass < array_size; pass++)
8	{
9	if (search_key == array[pass])





index	[0]	[1]	[2]	[3]	[4]	[5]	[6]	[7]	[8]	[9]	[10]	[11]
DATA	4	8	19	25	34	39	45	48	66	75	89	95

Figure 6.16 DATA

- a. The above figure shows an array **DATA** with the size of 12, which stores integer values. Determine the value of iteration, **pass**, **array[pass]**, **index** and number of comparisons in the following table format, in each steps of searching process on the array **DATA** by using the sequential search function in Program 6.11 for the following search keys.
  - i. 45
  - ii. 22

Iteration	pass	array[pass]	index	Number of
				comparisons
				••

vi.

Program 6.12 is a binary search function. Based on the same array figure above, answer all the following questions.

_	
1	// Program 6.12
2	int binary_search( int search_key, int array_size,
3	const int array [] ){
4	bool found = false;
5	int index = -1 //-1 means record not found
6	int MIDDLE,
7	LEFT = 0,
8	RIGHT = array_size-1;
9	
10	while (( LEFT<= RIGHT ) && (!found))
11	{ MIDDLE = (LEFT + RIGHT ) / 2;
12	// Get middle index
13	if (array[MIDDLE] == search_key)
14	{ index = MIDDLE;
15	found = true;
16	}
17	else if (array[MIDDLE] > search_key)
18	RIGHT= MIDDLE– 1;
19	// search is focused
20	<pre>// on the left side of list</pre>





21	else
22	LEFT = MIDDLE + 1
23	<pre>// search is focused</pre>
24	<pre>// on the right side of the list</pre>
25	} //end while
26	return index;
27	} //end function

b. Based on the above array, determine the value of iteration, LEFT, RIGHT, MIDDLE, array[MIDDLE], index and number of comparisons like in the table below, in each steps of searching process on the array DATA by using the binary search function in Program 6.12 for the following search keys.
i. 45
ii. 22

Iteration	LEFT	RIGHT MIDDLE		array[MIDDLE]	index	Number of
						comparisons

c. Based on your answers in questions (a) and (b), what can you conclude on the efficiency of both search algorithms? Explain in term of efficiency and number of comparisons.

# EXERCISE 7: SEARCHING TECHNIQUES AND THE ANALYSIS

Answer the following questions based on the sorted array named **marks** shown in Figure 6.17.

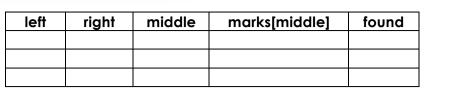
index	[0]	[1]	[2]	[3]	 [19]	[20]		[29]	[30]	 [34]	 [38]
marks	51	55	59	60	 75	76		85	86	 90	 95
							-				

a. Based on the given array in Figure 6.17, assume that your mark in this course is **60** and need to be searched using **SortedSeqSearch()** function. Show the tracing of the search using variables **index**, **p**, **search\_Key** and **found** as shown in the table format below.

index	search_Key	р	found

b. Assume your mark in this course is **90** and need to be searched using Binary Search function. Show the tracing of your search using variables **left**, **right**, **middle**, **marks[middle]** and **found** as shown in the table format below.





c. Fill in the following table with the number of steps and the complexity required in searching the **minimum mark**, **the average mark** (assume 75 is the average) and the **maximum mark**. Based on the results, compare and discuss the efficiency of Binary Search technique and Sequential Search (on sorted data) algorithms in the three searching cases.

	Search Comparisons				
Search Key	Linear Search		Binary Search		
	Number of steps	Complexity	Number of steps	Complexity	
51					
75					
95					
Efficiency Analysis for the three cases					